

# Kripke Completeness of Strictly Positive Modal Logics Over Meet Semilattices with Operators

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This talk is about a connection between various consequence relations for the fragment of propositional multi-modal logic that comprises implications  $\sigma \rightarrow \tau$ , where  $\sigma$  and  $\tau$  are *strictly positive modal formulas* constructed from propositional variables using conjunction  $\wedge$ , unary diamond operators  $\diamond_i$ , and the constant ‘truth’  $\top$ . We call such formulas *SP-implications* and we call sets of SP-implications *SP-theories*. These formulas can be interpreted on Kripke frames as well as on meet semilattices with unary monotone operators (SLOs). This gives rise to the following problem:

(**completeness**) identify SP-theories  $P$  that are *complete* in the sense that the two consequence relations  $P \models_{\mathbf{K}_r} \iota$  and  $P \models_{\text{SLO}} \iota$  coincide, where for any SP-implication  $\iota$ ,

$$\begin{aligned} P \models_{\mathbf{K}_r} \iota & \text{ iff } \iota \text{ is valid in every Kripke frame validating } P; \\ P \models_{\text{SLO}} \iota & \text{ iff } \iota \text{ is valid in every SLO validating } P. \end{aligned}$$

SP-implications are Sahlqvist, so for every modal formula  $\varphi$  and SP-theory  $P$

$$P \models_{\mathbf{K}_r} \varphi \text{ iff } \varphi \in \mathbf{K} + P \text{ iff } \varphi \approx \top \text{ is valid in every BAO validating } P, \quad (1)$$

where BAO stands for *Boolean algebras with normal and  $\vee$ -additive unary operators*. Note that, by (1), the completeness problem is equivalent to

(**conservativity**) the purely algebraic problem of whether the consequence relation  $P \models_{\text{BAO}}$  is *conservative* over  $P \models_{\text{SLO}}$  with respect to SP-implications, that is,  $P \models_{\text{SLO}} \iota$  iff  $P \models_{\text{BAO}} \iota$ , for any  $\iota$ ; and also to

(**axiomatisability**) the problem whether  $P$  *axiomatises* the SP-implicational fragment of the normal modal logic  $\mathbf{K} + P$  using the syntactic Birkhoff-type calculus corresponding to the algebraic consequence relation  $P \models_{\text{SLO}}$  (in other words, the problem whether  $P$  has a *modal companion*).

I am going to present two methods for proving completeness for SP-theories together with numerous sufficient conditions for their applicability. Note that incomplete SP-theories are easy to find, with two simplest ones being  $\{\diamond p \rightarrow p\}$  and  $\{\diamond p \rightarrow \diamond q\}$ .

This talk is based on a recent joint work with Agi Kurucz, Yoshihito Tanaka, Frank Wolter and Michael Zakharyashev accessible at <https://arxiv.org/pdf/1708.03403.pdf>